

# Logic Comprehensive Exam Summer 2018

## Three Parts

### Part Zero

1. State and prove the Tarski-Vaught criterion for elementary substructures.
2. State and prove Cantor's theorem on countable endless dense linear orders.
3. State the compactness theorem of first-order logic and sketch a proof of it.
4. State two equivalent formulations of the axiom of choice, and sketch a proof of their equivalence.

## Part One

**Do four of the following eight problems.**

1. Let  $\mathcal{L}$  be a first-order language and  $\mathfrak{M}$  be an  $\mathcal{L}$ -structure. Let  $\mathcal{L}(M)$  be  $\mathcal{L}$  expanded with constants for the individual elements of  $\mathfrak{M}$  and let  $\mathfrak{M}_M$  be the natural expansion of  $\mathfrak{M}$  to and  $\mathcal{L}(M)$ -structure. Show that for any  $\mathcal{L}$  structure  $\mathfrak{N}$  there is an elementary embedding  $f : \mathfrak{M} \rightarrow \mathfrak{N}$  if and only if there is an expansion of  $\mathfrak{N}$  to a model of  $\text{Th}(\mathfrak{M}_M)$ .
2. Let  $\mathcal{L}$  be the language with a single binary relation  $R$ . Let  $T$  be the theory axiomatized by stating that  $R$  is an equivalence relation with exactly two infinite classes and no other classes. Show that  $T$  has quantifier elimination.
3. Let  $\mathcal{L}$  be the first-order language consisting of countably many unary predicates  $U_i$  for  $i \in \omega$ . Let  $T$  be the  $\mathcal{L}$ -theory axiomatized by:
  - $\forall x \neg(U_i x \wedge U_j x)$  for all  $i \neq j \in \omega$ ;
  - $\exists x_1, \dots, x_n (\bigwedge_{1 \leq i \leq n} U_j x_i \wedge \bigwedge_{1 \leq i \neq k \leq n} x_i \neq x_k)$  for all  $j, n \in \omega$ .

Describe all of the countable models of  $T$  and indicate which, if any, are prime or saturated. Justify your answer.

4. Let  $\mathfrak{M}_i$  for  $i \in \omega$  be finite  $\mathcal{L}$ -structures. Let  $\mathcal{U}$  be an ultrafilter on  $\omega$ . Show that in the ultraproduct  $\mathfrak{M} = (\prod_{i \in \omega} \mathfrak{M}_i) / \mathcal{U}$  there is no formula  $\varphi(x, y)$  which defines an infinite linear order with no largest element.
5. Let  $\mathcal{L}$  be a first order language consisting of a binary relation symbol  $<$  and a unary predicate  $P$ . Let  $\mathfrak{M}$  be the  $\mathcal{L}$ -structure with universe the real numbers where  $<$  is interpreted as the usual order on the real numbers and  $P$  is interpreted as the set of rational numbers. Show that  $\text{Th}(\mathfrak{M})$  is  $\omega$ -categorical.
6. Let  $\mathcal{L}$  be a countable first-order language. Show that a countable  $\mathcal{L}$ -structure  $\mathfrak{M}$  is countably saturated if and only if it is countably universal and homogeneous.
7. Let  $T$  be a theory in a countable language  $\mathcal{L}$ . Let  $\varphi_i(x)$  for  $i \in \omega$  be a collection of formulas in the free variable  $x$  and let  $\psi(x)$  be another such formula. Suppose that for any other formula  $\sigma(x)$  if  $T \models \forall x(\sigma(x) \rightarrow \psi(x))$  then  $T \models \exists x(\sigma(x) \wedge \varphi_i(x))$  for some  $i \in \omega$ . Show that there is a model  $\mathfrak{M}$  of  $T$  so that  $\psi(\mathfrak{M}) \subseteq \bigcup_{i \in \omega} \varphi_i(\mathfrak{M})$ .
8. A theory  $T$  without finite models is called *strongly minimal* if in any model  $\mathfrak{M}$  of  $T$  any definable subset  $X \subseteq M$  is either finite or co-finite. (More specifically if  $\mathfrak{M} \models T$ ,  $\varphi(x, \bar{y})$  is a formula, and  $\bar{a} \in M^{|\bar{y}|}$  then the set of realizations of  $\varphi(x, \bar{a})$  in  $\mathfrak{M}$  is either finite or co-finite in  $M$ .) Suppose that  $T$  is strongly minimal and show that if  $\varphi(x, \bar{y})$  is any formula then there is a natural number  $N$  so that if  $\mathfrak{M} \models T$  and  $\bar{a} \in M^{|\bar{y}|}$  then if  $|\{b \in M : \mathfrak{M} \models \varphi(b, \bar{a})\}| > N$  then  $\{b \in M : \mathfrak{M} \models \varphi(b, \bar{a})\}$  is infinite.

## Part Two

**Do four of the following eight problems.**

1. Let  $T$  be a recursive, consistent set of sentences containing the axioms of ZFC. A sentence  $\varphi$  is *independent* of  $T$  if neither  $\varphi$  nor  $\neg\varphi$  are provable from  $T$ . Show that the set of sentences that are independent of  $T$  is not recursively enumerable.
2. Show that a language is decidable iff it has an enumerator with the property that if it outputs a word  $v$  before a word  $w$ , then the length of  $v$  is less than or equal to the length of  $w$ .
3. Let  $\{\ulcorner M_1 \urcorner, \ulcorner M_2 \urcorner, \dots\}$  be a recursively enumerable language consisting of descriptions of deciders  $M_1, M_2, \dots$ . Suppose the underlying alphabet consists of the symbols 0 and 1. Show that there is a decider whose language is different from the languages decided by  $M_1, M_2, \dots$ .
4. Show in ZF, assuming the consistency of Peano arithmetic, that there is a bijection between  $\mathcal{P}(\omega)$  and the set of complete and consistent theories (in the language of arithmetic) extending Peano Arithmetic.
5. Let  $\kappa$  be an uncountable regular cardinal, and suppose that  $\langle X_\alpha \mid \alpha < \theta \rangle$  is a strictly increasing sequence (of length  $\theta$ ) of subsets of  $\kappa$ , such that  $\bigcup_{\alpha < \theta} X_\alpha = \kappa$ . Show:
  - (a)  $\theta < \kappa^+$ .
  - (b) for every nonzero  $\gamma < \kappa^+$ , there is such a sequence of length  $\gamma$ .
  - (c) Assuming that for each  $\alpha < \theta$ , the cardinality of  $X_\alpha$  is less than  $\kappa$ , it follows that  $\theta = \kappa$ .
6. The principle of Dependent Choices (DC) says that whenever  $\langle u, r \rangle$  is a binary system with the property that for every  $y \in u$ , there is an  $x \in u$  such that  $xry$ , it follows that there is an infinite decreasing chain in  $\langle u, r \rangle$ , that is, a function  $f : \omega \rightarrow u$  such that for all  $n \in \omega$ ,  $f(n+1)rf(n)$ . The countable Löwenheim-Skolem Theorem ( $\text{LS}_\omega$ ) says that every infinite model  $M$  of a countable first order language has a countable elementary submodel. Show in ZF that  $\text{LS}_\omega$  is equivalent to DC.
7. Let  $\kappa$  be a strongly inaccessible cardinal, and let  $M = \langle V_\kappa, \in \upharpoonright V_\kappa \rangle$ , a model of the language of set theory. Let  $N = \langle X, \in \upharpoonright X \rangle \prec M$  be a countable elementary submodel.
  - (a) Show that  $N \cap \omega_1 \in \omega_1$  and  $\omega_1 \in X$ .
  - (b) Show that  $X$  is not transitive.
  - (c) Show that  $N$  is extensional and well-founded.
  - (d) Let  $f : N \rightarrow \langle U, \in \upharpoonright U \rangle$  be the Mostowski collapse and isomorphism. Show that  $f \upharpoonright \omega_1$  is the identity and  $f(\omega_1) = \omega_1 \cap X$ .
8. Prove the following statements (in ZFC):
  - (a) Let  $u$  be a transitive set. Whenever  $\varphi(\vec{x})$  is a  $\Delta_0$ -formula and  $\vec{a} \in u$ , then
 
$$\varphi(\vec{a}) \text{ holds iff } \langle u, \in \upharpoonright u \rangle \models \varphi(\vec{a}).$$
  - (b) Let  $\psi(\vec{x})$  be a  $\Sigma_1$ -formula, let  $\kappa$  be an infinite cardinal, let  $\vec{a} \in H_\kappa$  (where  $H_\kappa = \{b \mid \text{card}(b) < \kappa\}$ ). Then
 
$$\psi(\vec{a}) \text{ holds iff } \langle H_\kappa, \in \upharpoonright H_\kappa \rangle \models \psi(\vec{a}).$$